Network flow algorithms

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Introduction

- Input: a directed graph where each edge is a pipe with some capacity
- How much flow you can go from the source to the sink?



Concepts



Figure 2b - The residual network of the network in 2a

Figure 3a

residual network

Ford-Fulkerson algorithm

 Keep finding augmenting paths in the residual network while they exist



Figure 1a - Maximum Flow in a network

Figure 1b - The residual network of the network in 1a

Ford-Fulkerson

```
def max_flow(graph, source, sink):
flow = 0
while capacity := graph.find_augmenting_path(source, sink):
    flow += capacity
    return flow
```

Edmonds-Karp

- Use BFS to find augmenting paths!
- $O(VE^2)$



Dinitz's algorithm

- Use BFS to find augmenting paths!
- $O(V^2 E)$ vs. $O(VE^2)_{E-k}$



Max-flow min-cut

 Maximum flow solves another problem: min-cut in a graph (minimal set of edges that need to be removed to make a node unreachable from another node)



The minimum cut of this network is the sum of the capacities of XB and CY and equals 3, which is also the value of the maximum flow.

Figure 5 - A minimum cut in a network



• Multi-source multi-sink



Figure 6 - Reduction of a multiple-source / multiple-sink max-flow problem

Maximum bipartite matchings



Modeling: Titanic problem



Min-cost max-flow

- Sometimes edges have weights as well, and we want to pick a matching that has the min/max total weight
- Assignment problem: each person has job preference & salary expectations



Problems, problems

- Network flow:
 - UVA 00259, 00820, 00563, 11167, 11380, 12873
 - Kattis: unfairplay
 - Extra:
 - CP4.2: page 434
 - UVa: 00753, 10122, 10330, 10511, 10735
 - Topcoder: rookattack, graduation, parking
- Matching:
 - UVA 11138, 00670, 12644
 - [mcmf] UVA 10594, 11301, 10746